

# 51

"Wisdom is better than weapons of war."

## Decompilation / EXE to C

Decompilation is the reverse of compilation. That is, we can get a C file from EXE file! The most important problem in converting back C file from EXE file is loss of variable names and loss of function names. Machine code won't store variable names. So it is not at all possible to get back the original C code.

### 51.1 Basic Idea

Since it is a reverse of compilation, we must analyze how a compiler works and the corresponding machine code for the functions like `printf( )`, `scanf( )` etc. In other words, we must find the 'signature' of each C functions and C statements.

### 51.2 DCC

#### 51.2.1 Disclaimer

DCC is a decompiler written by **Cristina Cifuentes** and **Mike Van Emmerik** while at the Queensland University of Technology, Australia. Copyright is owned by **Cristina Cifuentes** and the Queensland University of Technology. DCC is merely a prototype tool and more work needs to be done in order to have a fully working decompiler.

#### Important Notice

I have received permission to use the article about DCC from the authors (Cristina Cifuentes and Mike Van Emmerik) with the condition of including the above disclaimer note. As Cristina Cifuentes and Mike Van Emmerik are not currently involving in decompilation, it seems they don't like to receive any request or correspondence regarding their decompilation work. So the reader is requested **not** to disturb them.

#### 51.2.2 Notice

Decompilation is a technique that allows you to recover lost source code. It is also needed in some cases for computer security, interoperability and error correction. dcc, and any decompiler in general, should not be used for "cracking" other programs, as programs are protected by copyright. Cracking of programs is not only illegal but it rides on other's creative effort.

### 51.2.3 DCC Facts

The dcc decompiler decompiles .exe files from the (i386, DOS) platform to C programs. The final C program contains assembler code for any subroutines that are not possible to be decompiled at a higher level than assembler.

The analysis performed by dcc is based on traditional compiler optimization techniques and graph theory. The former is capable of eliminating registers and intermediate instructions to reconstruct high-level statements; the later is capable of determining the control structures in each subroutine.

Please note that at present, only C source is produced; dcc cannot (as yet) produce C++ source.

The structure of a decompiler resembles that of a compiler: a front-, middle-, and back-end which perform separate tasks. The front-end is a machine-language dependent module that reads in machine code for a particular machine and transforms it into an intermediate, machine-independent representation of the program. The middle-end (aka the Universal Decompiling Machine or UDM) is a machine and language independent module that performs the core of the decompiling analysis: data flow and control flow analysis. Finally, the back-end is high-level language dependent and generates code for the program (C in the case of dcc).

In practice, several programs are used with the decompiler to create the high-level program. These programs aid in the detection of compiler and library signatures, hence augmenting the readability of programs and eliminating compiler start-up and library routines from the decompilation analysis.

### 51.2.4 Example of Decompilation

We illustrate the decompilation of a fibonacci program (see Figure 4). Figure 1 illustrates the relevant machine code of this binary. No library or compiler start up code is included. Figure 2 presents the disassembly of the binary program. All calls to library routines were detected by dccSign (the signature matcher), and thus not included in the analysis. Figure 3 is the final output from dcc. This C program can be compared with the original C program in Figure 4.

```

55 8B EC 83 EC 04 56 57 1E B8 94 00 50 9A
0E 00 3C 17 59 59 16 8D 46 FC 50 1E B8 B1 00 50
9A 07 00 F0 17 83 C4 08 BE 01 00 EB 3B 1E B8 B4
00 50 9A 0E 00 3C 17 59 59 16 8D 46 FE 50 1E B8
C3 00 50 9A 07 00 F0 17 83 C4 08 FF 76 FE 9A 7C
00 3B 16 59 8B F8 57 FF 76 FE 1E B8 C6 00 50 9A
0E 00 3C 17 83 C4 08 46 3B 76 FC 7E C0 33 C0 50
9A 0A 00 49 16 59 5F 5E 8B E5 5D CB 55 8B EC 56
8B 76 06 83 FE 02 7E 1E 8B C6 48 50 0E E8 EC FF
59 50 8B C6 05 FE FF 50 0E E8 E0 FF 59 8B D0 58
03 C2 EB 07 EB 05 B8 01 00 EB 00 5E 5D CB

```

**Figure 1 - Machine Code for Fibonacci.exe**

```

proc_1 PROC FAR
000 00053C 55          PUSH          bp
001 00053D 8BEC          MOV          bp, sp
002 00053F 56          PUSH          si
003 000540 8B7606       MOV          si, [bp+6]
004 000543 83FE02       CMP          si, 2
005 000546 7E1E          JLE         L1
006 000548 8BC6          MOV          ax, si
007 00054A 48          DEC          ax
008 00054B 50          PUSH          ax
009 00054C 0E          PUSH          cs
010 00054D E8ECFF       CALL        near ptr proc_1
011 000550 59          POP          cx
012 000551 50          PUSH          ax
013 000552 8BC6          MOV          ax, si
014 000554 05FEFF       ADD          ax, 0FFFEh
015 000557 50          PUSH          ax
016 000558 0E          PUSH          cs
017 000559 E8E0FF       CALL        near ptr proc_1
018 00055C 59          POP          cx
019 00055D 8BD0          MOV          dx, ax
020 00055F 58          POP          ax
021 000560 03C2          ADD          ax, dx
022 000561 5E          POP          si
023 00056B 5E          L2: POP          si
024 00056C 5D          POP          bp
025 00056D CB          RETF
026 000566 B80100       L1: MOV          ax, 1
027 000569 EB00          JMP         L2
proc_1 ENDP

```

```

main PROC FAR
000 0004C2 55          PUSH          bp
001 0004C3 8BEC          MOV          bp, sp
002 0004C5 83EC04       SUB          sp, 4
003 0004C8 56          PUSH          si
004 0004C9 57          PUSH          di
005 0004CA 1E          PUSH          ds
006 0004CB B89400       MOV          ax, 94h
007 0004CE 50          PUSH          ax
008 0004CF 9A0E004D01  CALL        far ptr printf
009 0004D4 59          POP          cx
010 0004D5 59          POP          cx
011 0004D6 16          PUSH          ss
012 0004D7 8D46FC       LEA         ax, [bp-4]
013 0004DA 50          PUSH          ax
014 0004DB 1E          PUSH          ds
015 0004DC B8B100       MOV          ax, 0B1h

```

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```
016 0004DF 50          PUSH      ax
017 0004E0 9A07000102  CALL     far ptr scanf
018 0004E5 83C408          ADD      sp, 8
019 0004E8 BE0100          MOV      si, 1
021 000528 3B76FC          L3:     CMP      si, [bp-4]
022 00052B 7EC0          JLE     L4
023 00052D 33C0          XOR      ax, ax
024 00052F 50          PUSH     ax
025 000530 9A0A005A00  CALL     far ptr exit
026 000535 59          POP      cx
027 000536 5F          POP      di
028 000537 5E          POP      si
029 000538 8BE5          MOV      sp, bp
030 00053A 5D          POP      bp
031 00053B CB          RETF
032 0004ED 1E          L4:     PUSH     ds
033 0004EE B8B400          MOV      ax, 0B4h
034 0004F1 50          PUSH     ax
035 0004F2 9A0E004D01  CALL     far ptr printf
036 0004F7 59          POP      cx
037 0004F8 59          POP      cx
038 0004F9 16          PUSH     ss
039 0004FA 8D46FE          LEA     ax, [bp-2]
040 0004FD 50          PUSH     ax
041 0004FE 1E          PUSH     ds
042 0004FF B8C300          MOV      ax, 0C3h
043 000502 50          PUSH     ax
044 000503 9A07000102  CALL     far ptr scanf
045 000508 83C408          ADD      sp, 8
046 00050B FF76FE          PUSH     word ptr [bp-2]
047 00050E 9A7C004C00  CALL     far ptr proc_1
048 000513 59          POP      cx
049 000514 8BF8          MOV      di, ax
050 000516 57          PUSH     di
051 000517 FF76FE          PUSH     word ptr [bp-2]
052 00051A 1E          PUSH     ds
053 00051B B8C600          MOV      ax, 0C6h
054 00051E 50          PUSH     ax
055 00051F 9A0E004D01  CALL     far ptr printf
056 000524 83C408          ADD      sp, 8
057 000527 46          INC      si
058          JMP      L3          ;Synthetic inst

      main ENDP
```

**Figure 2 - Code produced by the Disassembler**

```

/*
 * Input file   : fibo.exe
 * File type    : EXE
 */
int proc_1 (int arg0)
/* Takes 2 bytes of parameters.
 * High-level language prologue code.
 * C calling convention.
 */
{
int loc1;
int loc2; /* ax */

    loc1 = arg0;
    if (loc1 > 2) {
        loc2 = (proc_1 ((loc1 - 1)) + proc_1 ((loc1 + 0xFFFFE)));
    }
    else {
        loc2 = 1;
    }
    return (loc2);
}

void main ( )
/* Takes no parameters.
 * High-level language prologue code.
 */
{
int loc1;
int loc2;
int loc3;
int loc4;

    printf ("Input number of iterations: ");
    scanf ("%d", &loc1);
    loc3 = 1;
    while ((loc3 <= loc1)) {
        printf ("Input number: ");
        scanf ("%d", &loc2);
        loc4 = proc_1 (loc2);
        printf ("fibonacci(%d) = %u\n", loc2, loc4);
        loc3 = (loc3 + 1);
    } /* end of while */
    exit (0);
}

```

**Figure 3 - Code produced by dcc in C**

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```
#include <stdio.h>

int main( )
{ int i, numtimes, number;
  unsigned value, fib();

  printf("Input number of iterations: ");
  scanf ("%d", &numtimes);
  for (i = 1; i <= numtimes; i++)
  {
    printf ("Input number: ");
    scanf ("%d", &number);
    value = fib(number);
    printf("fibonacci(%d) = %u\n", number, value);
  }
  exit(0);
}

unsigned fib(x)          /* compute fibonacci number recursively */
int x;
{
  if (x > 2)
    return (fib(x - 1) + fib(x - 2));
  else
    return (1);
}
```

**Figure 4 – Initial / Original C Program**